

Reinforcement Learning for 5G Caching with Dynamic Cost

*A. Sadeghi, F. Sheikholeslami, **A. G. Marques** and G. B. Giannakis*

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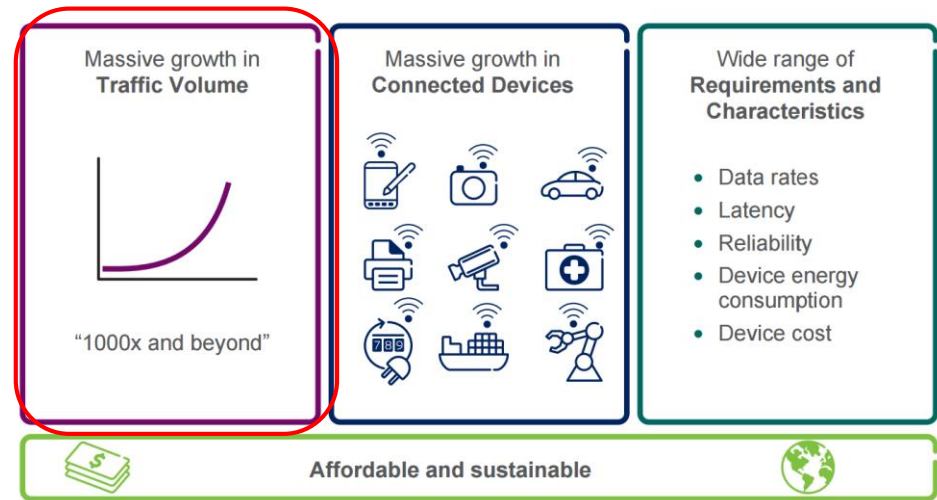


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Challenges for 5G

❑ Evolving user profiles

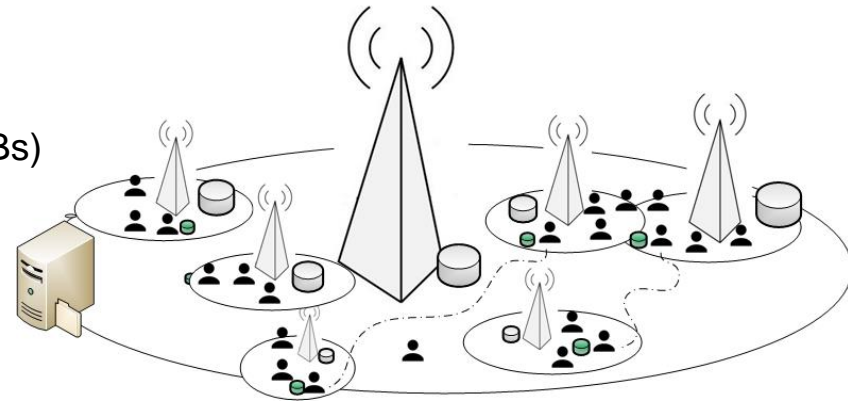
- Mobile video streaming
 - ✓ 50 % of traffic volume in 4G
 - ✓ 500-fold increase in 5G
- Social networking
 - ✓ 15% of traffic volume in 4G
- Music, games ...



❑ 60% of data is reusable, a.k.a. **contents**

❑ Heterogeneous network architecture (HetNet)

- Utilization of **storage** units at small base stations (SBs)
 - Proactively store popular contents (**cache**)
- Challenge: **what** and **when** to store?
 - Requires **learning** content popularities



Caching in wireless networks

❑ Memory-enabled SBs

- Cache during off-peak hours
- Reduce load on backhauls during peak traffic periods
- Reduce cost for providing service with high QoS

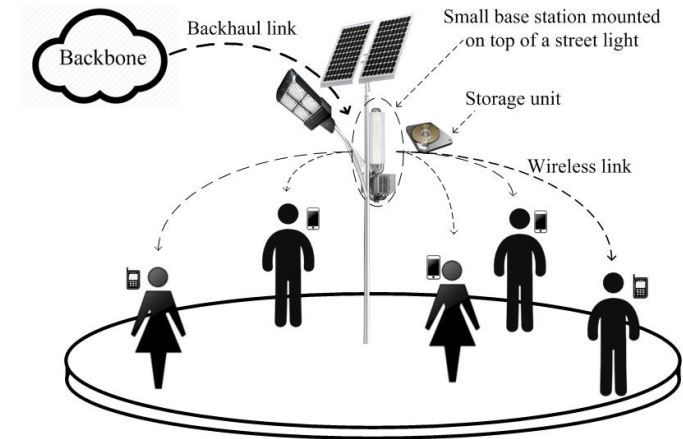
❑ Generally unknown content popularity profiles

❑ Prior art

- Popularity profile learning [Blasco, Gunduz'15], [Baştuğ', Debbah, Saad'16], [Bharath'16]
- Multi-armed bandit (MAB) formulation [Belasco et al'14]
- Distributed and convexified MAB [Sengupta et al'14]
- Dynamic popularity profiles [Sadeghi et al'18]
- Game-theoretic caching [Hamidouche, Debbah, Saad'16]
- Coded caching [Maddah-Ali, Usir'16], [Alizadeh, Avestimehr'16], [Amiri, Gunduz'17]

⇒ Unknown static popularity profile

⇒ Unknown dynamic popularity profile



Proposed approach

Caching via reinforcement learning while considering dynamic fetching-caching costs

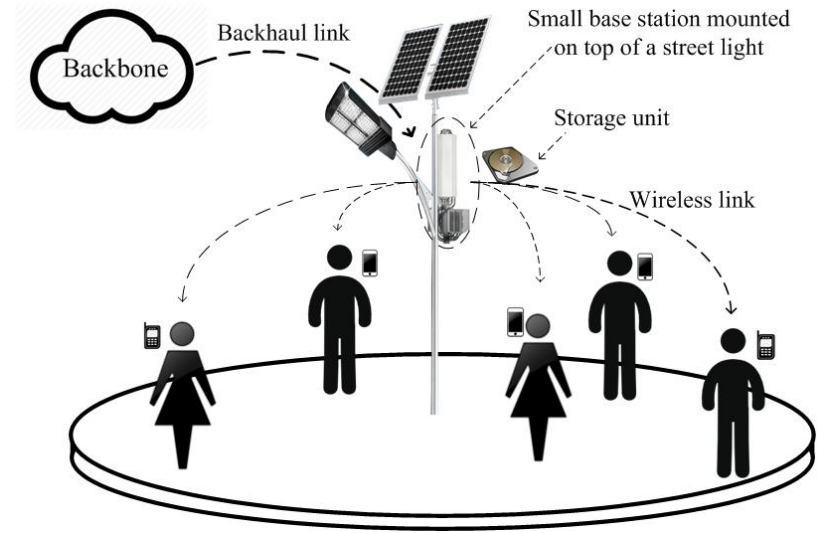
Problem statement

❑ Discrete-time network

- Access point with storage to cache popular files
- Total number of F contents at the back-bone
- User content requests are served using
 - ✓ Proactively cached contents, or
 - ✓ Reactively fetched (via back-haul link) contents

❑ Pertinent costs for provisioning a request

- Storing locally in access point: processing, storage and energy consumption
- Fetching from cloud: scheduling, routing, transmission through expensive back-haul



Minimize sum-average cost by sequentially making fetching-caching decisions

Content-dependent variable costs

□ File- and size-dependent caching-fetching costs

- Cost of caching file f at slot t

$$c_t^f := \overset{\text{size dep.}}{\sigma_f}(c'_t + c'^f_t) + \overset{\text{file dep.}}{(c''_t + c''^f_t)}$$

- Cost of fetching file f at slot t

$$\phi_t^f := \sigma_f(\phi'_t + \phi'^f_t) + (\phi''_t + \phi''^f_t)$$

□ Base-station receives user file request for file f at slot t ($r_t^f = 1$)

- Fetching-action variable $w_t^f \in \{0, 1\}$
- Caching-action variable $a_t^f \in \{0, 1\}$

Goal: Given pdf of iid $\{r_\tau^f, c_\tau^f, \phi_\tau^f\}_{\tau=t+1}^\infty$ and instantaneous values per slot t , find $\{w_\tau^{f*}, a_\tau^{f*}\}_{\tau=t+1}^\infty$

Fetch and cache via dynamic programming

- ❑ Cache-state variable $s_t^f \in \{0, 1\}$ where $s_t^f = a_{t-1}^f$
- ❑ Constraints on fetch-cache decision variables $\{w_\tau^f, a_\tau^f\}_{\tau=t+1}^\infty$
 - File requests must be served (no drop-off allowed) $r_t^f \leq w_t^f + s_t^f, \quad \forall f, t$
 - Caching is feasible iff file is available $a_t^f \leq s_t^f + w_t^f, \quad \forall f, t$
- ❑ Fetch-cache cost $C^f(a_t^f, w_t^f; c_t^f, \phi_t^f) = c_t^f a_t^f + \phi_t^f w_t^f$


$$\begin{aligned} \min_{\{(w_\tau^f, a_\tau^f)\}_{f, \tau \geq t}} \quad & \bar{C}_t := \sum_{\tau=t}^{\infty} \sum_{f=1}^F \gamma^{\tau-t} \mathbb{E} [C^f(a_\tau^f, w_\tau^f; c_\tau^f, \phi_\tau^f)] \\ \text{s.t.} \quad & (w_\tau^f, a_\tau^f) \in \mathcal{X}(r_\tau^f, s_\tau^f), \quad \forall f, \tau \geq t \end{aligned}$$

$$\mathcal{X}(r_t^f, s_t^f) := \{(w, a) \mid w, a \in \{0, 1\}, s_t^f = a_{t-1}^f, r_t^f \leq w + s_t^f, a \leq s_t^f + w\}$$

- ❑ Optimization separable across files!

Formulating per content optimization

□ Given $(s_t^f, r_t^f, c_t^f, \phi_t^f)$

$$\theta_t^f := [r_t^f, c_t^f, \phi_t^f]$$


$$\begin{aligned} (w_t^{f*}, a_t^{f*}) &:= \arg \min_{(w,a) \in \mathcal{X}(r_t^f, s_t^f)} \left\{ \mathbb{E}_{r_\tau^f, c_\tau^f, \phi_\tau^f; \tau > t} \left[\min_{(w_\tau, a_\tau) \in \mathcal{X}(r_\tau^f, s_\tau^f)} \left\{ \sum_{\tau=t}^{\infty} \gamma^{\tau-t} \left[C^f(a_\tau^f, w_\tau^f; c_\tau^f, \phi_\tau^f) \mid a_t^f = a, w_t^f = w, \theta_t^f = \theta_0^f \right] \right\} \right] \right\} \\ &= \arg \min_{(w,a) \in \mathcal{X}(r_t^f, s_t^f)} \left\{ c^f(a, w; c_t^f, \phi_t^f) + \mathbb{E}_{r_\tau^f, c_\tau^f, \phi_\tau^f; \tau > t} \left[\min_{(w_\tau, a_\tau) \in \mathcal{X}(r_\tau^f, s_\tau^f)} \sum_{\tau=t+1}^{\infty} \gamma^{\tau-t} \left[c^f(a_\tau^f, w_\tau^f; c_\tau^f, \phi_\tau^f) \mid s_{t+1}^f = a \right] \right] \right\} \end{aligned}$$

□ Marginalized value function

$$\bar{V}^f(s^f) := \mathbb{E}_{r^f, c^f, \phi^f} \left[\min_{(w,a) \in \mathcal{X}(r^f, s^f)} \{ C^f(a, w; c^f, \phi^f) + \gamma \bar{V}^f(a) \} \right]$$

Value iteration algorithm

□ Value iteration algorithm to find $\bar{V}^f(s^f)$

- **Input:** probability density functions of c^f, ϕ^f and r^f
- **Initialize:** $\bar{V}_0^f(s) = 0$, for $s \in \{0, 1\}$
- **While** $|\bar{V}_k^f(s) - \bar{V}_{k+1}^f(s)| < \epsilon; \quad \forall f, s \in \{0, 1\}$

$$\bar{V}_{k+1}^f(s) = \mathbb{E}_{r^f, c^f, \phi^f} \min_{(w, a) \in \mathcal{X}(r^f, s)} \left\{ C^f(a, w; c^f, \phi^f) + \gamma \bar{V}_k^f(a) \right\} \quad \text{for } s = 0, 1$$

- **Output:** $\bar{V}^f(0), \bar{V}^f(1)$

□ Optimal cache-fetch decisions given $(s_t^f, r_t^f, c_t^f, \phi_t^f)$

$$(w_t^{f*}, a_t^{f*}) = \arg \min_{(w, a) \in \mathcal{X}(r_t^f, s_t^f)} \left\{ C^f(a, w; c_t^f, \phi_t^f) + \bar{V}^f(a) \right\}$$

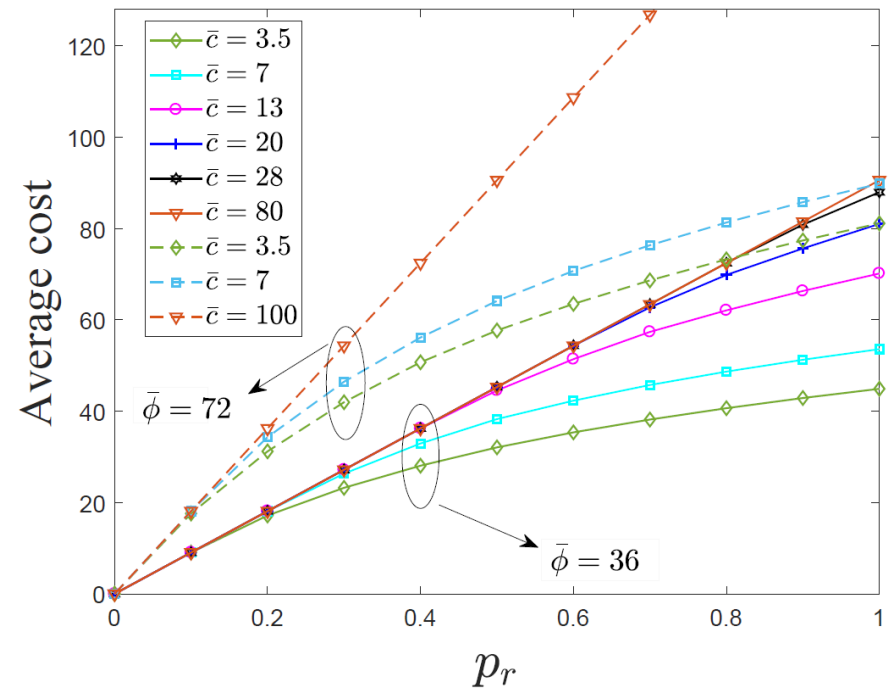
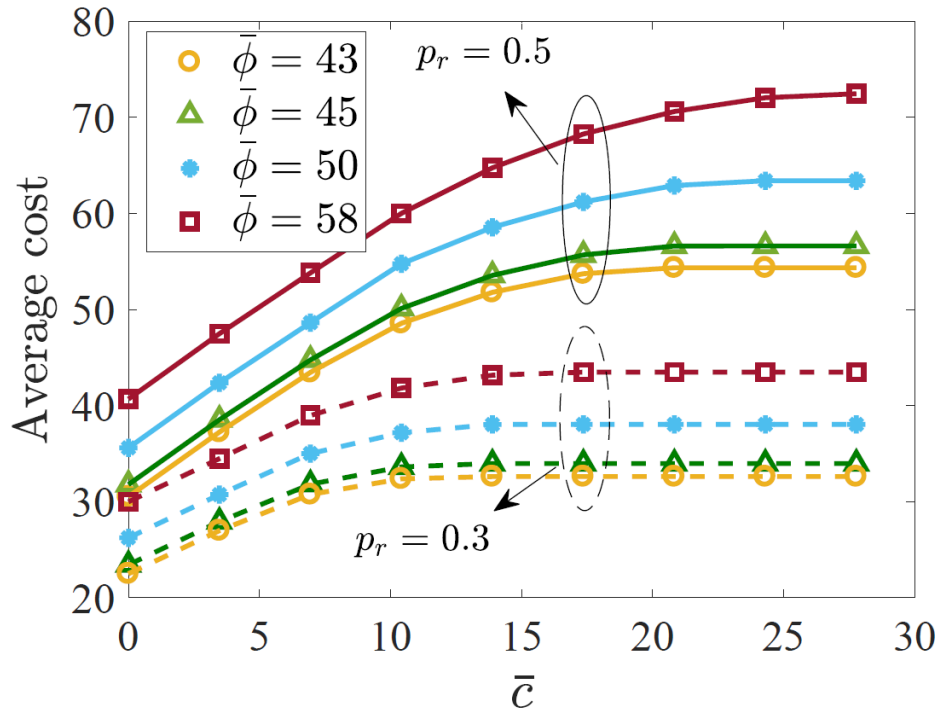
□ Under iid assumption, marginalized value function takes binary input

\Rightarrow Fast convergence!

Numerical test with dynamic cost

- Per file f pdfs

$$c_t^f \sim \mathcal{U}(0, 2\bar{c}), \quad \phi_t^f \sim \mathcal{U}(0, 2\bar{\phi}), \quad r_t^f \sim \text{Bernoulli}(p_r^f)$$

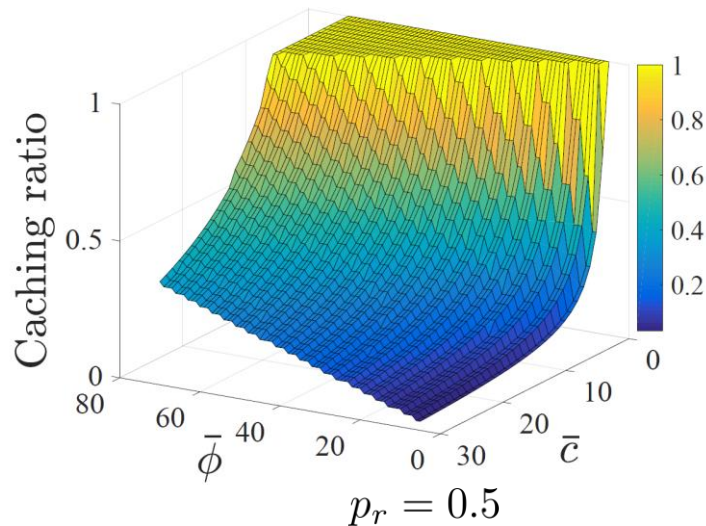


- ✓ For fixed $\bar{\phi}$, up to a certain value of \bar{c} , caching is encouraged
- ✓ Large p_r encourages caching for a larger range of \bar{c}
- ✓ Increased cost with more requests p_r

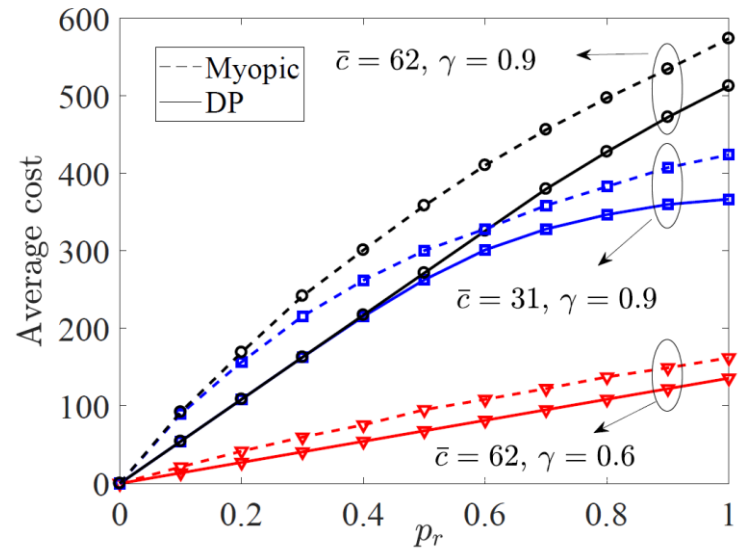
Further comparisons

- $c_t^f \sim \mathcal{U}(0, 2\bar{c})$, $\phi_t^f \sim \mathcal{U}(0, 2\bar{\phi})$, $r_t^f \sim \text{Bernoulli}(p_r^f)$

caching ratio := $\frac{\text{No. of caching decisions}}{\text{Total No. of decisions}}$



✓ DP vs. myopic caching



- ✓ Low \bar{c} , high $\bar{\phi}$ \longrightarrow caching ratio = 1 (flat area)
- ✓ High \bar{c} , low $\bar{\phi}$ \longrightarrow caching ratio = 0
- ✓ Intermediate values \longrightarrow $0 < \text{caching ratio} < 1$
- ✓ DP considers future, thus reaches smaller average cost versus myopic caching

Future research and stakeholder analysis

- ❑ Multi-file caching considering queuing and cache refreshing costs
- ❑ Cooperative caching across neighboring small cells
- ❑ Cross-layer design of coded caching
- ❑ Privacy-preserving, secure, space-time variable caching

Thank you!